# Silo/HDF5 Modifications for Dawn

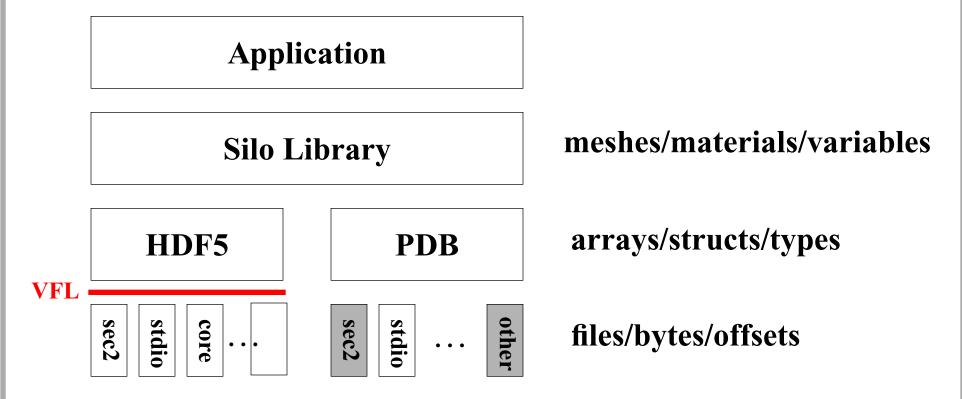
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# Silo Background



# **Benefits (= flexibility)**

- platform independent, self-describing, archiveable data
- random access (more true of post-processors than simulation codes)

# **Drawbacks (= performance degradation)**

- metadata (data a lib writes on behalf of its caller)
- caller is far removed from actual disk I/O behavior/control

# Poor Man's Parallel I/O

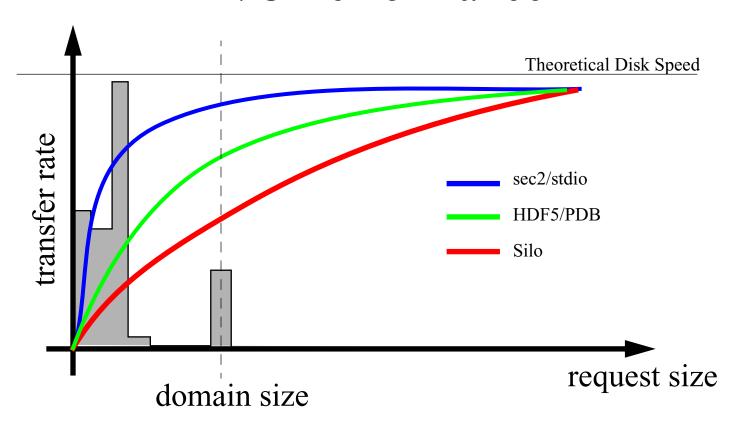
# Truly concurrent, parallel I/O to a single file is problematic

- Difficult to make perform well even for relatively simple I/O patterns.
- The global monolithic "whole" object is decomposed on read, re-composed on write
- Does not support multi-physics codes where I/O patterns are more complex

## Poor Man's Parallel I/O: Parallelism at the price of multiple files

- Serial I/O to multiple files, simultaneously
- #files != #MPI-tasks
- Very flexible with what each MPI-task needs to do in the way of I/O
- Do not pay cost of "decomposing on read" and "recomposing on write"
- Note: Lustre can't tell the difference (almost)

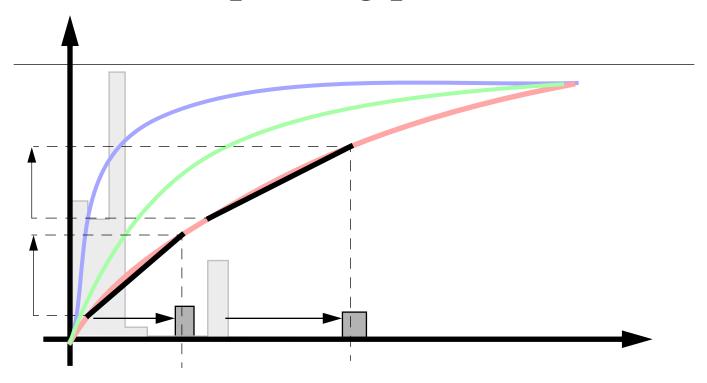
# I/O Performance



# Histogram of a recent Ares dump

	writes	bytes	%writes	cum.%writes	%bytes
<10^1 bytes:	48	217	20.1680	20.1680	.0001
<10^2 bytes:	41	1485	17.2268	37.3949	.0009
<10^3 bytes:	116	22474	48.7394	86.1344	.0136
<10^4 bytes:	8	30540	3.3613	89.4957	.0186
<10 <sup>5</sup> bytes:	0	0	0	89.4957	0
<10^6 bytes:	3	1092492	1.2605	90.7563	.6655
<10^7 bytes:	22	162989412	9.2436	100.0000	99.3010

# Aggregation is key to improving performance



# Aggregation

- Gather many smaller requests into fewer larger ones
- Need memory to do this.
- Try aggregating as much as possible WITHIN one MPI-task first.
- Failing that, start aggregating ACROSS MPI-tasks.

# Simple Aggregation Strategies

#### **HDF5's Core VFD:**

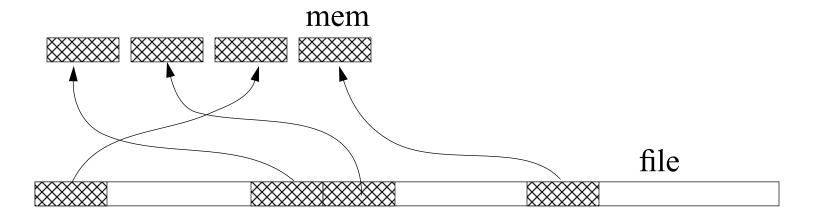
- Stores everything to a growing buffer in memory.
- Writes buffer to file on close.
- Reads ENTIRE file to memory buffer on open.
- Represents upper-bound of what is possible at expense of (a lot) of memory.
- Only works if when code does I/O, it is dumping less than 50% of available memory.
- Not a good long term solution

## HDF5's Split VFD:

- Splits data into two classes; raw and meta, writing each to its own VFD.
- Metadata uses core VFD, raw data uses sec2 VFD
- Improves performance but at price of two files on disk per one created by app.

# Silo's new Block VFD for Dawn

#### Breaks virtual file into blocks



## Does I/O only in blocks

• Allocates enough memory to keep N blocks in memory; uses LRU to pre-empt.

# Two Parameters set by code

- SILO\_BLOCK\_SIZE ( should be multiple of filesystem blocksize)
- SILO\_BLOCK\_COUNT (more is better)

#### **Good Values for Dawn**

- *SILO\_BLOCK\_SIZE* = (1<<20)
- SILO\_BLOCK\_COUNT=16 (16 Megabytes total)

# Other VFDs We May Write

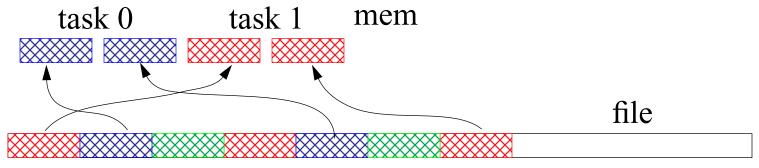
#### **Remote-Core VFD**

- Use extra MPI-tasks just for I/O
- Code "writes" to memory in these extra MPI-tasks through enhanced core VFD
- Code goes back to compute while data drains to files from the extra MPI-tasks
- Should be absolute fastest as code doesn't ever wait for disk; just MPI-send(s).

### **Smart-Split VFD:**

- Only one file is produced
- Raw data is block buffered as in new Silo VFD
- Metadata is kept in memory until file close, then tacked onto end of file

# **Extend Block VFD to stripe across MPI-tasks**



- Let application "think" its writing to different files
- What if each MPI-task is writing wildly different amounts of data?
- May be possible to make this completely transparent to HDF5